

A mathematician's journey with Grothendieck

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Mathematics, history, literature and Latin

- At school and high-school,
mathematics was an easy subject
(later, it became more challenging)
but I had other passions.
- Early passions:
 - History (already at elementary school age).
 - Literature (a little later):
in particular French and Russian literature.
 - Latin language (especially at age 13-15).

Remarks:

- Mathematics has a very strong historical dimension.
It is a collective tradition,
built century after century, even millenia.
- In some sense, mathematics is a part of literature.
- Ancien Latin is a very structural language.
Studying Latin was a good preparation for math. logic.

First discovery of Grothendieck's mathematics

- I think I had never heard about Grothendieck before my second year at “Ecole Normale Supérieure”, at Master level.
- I discovered Grothendieck because an “initiation to algebraic geometry” seminar had been organized for students. I wondered what “algebraic geometry” could mean, and I decided to participate in this seminar.
- I quickly realized that algebraic geometry was a very old domain which had been given entirely new very modern foundations and methods by essentially one mathematician.
- Starting from this discovery moment, I spent several years studying EGA and SGA ...

Scheme theory

Difficult to say a few words about a continent . . .

Two fascinating features of this continent were:

- The very definition of geometric objects (schemes) as topological spaces (of a strange kind: they were not Hausdorff!) endowed with a “sheaf of commutative rings” carrying the geometric structure.

→ This idea came from Jean-Pierre Serre and had been lifted by Grothendieck to the natural level of generality.

- Schemes were not considered individually but through their relations with other schemes, formalized as a “category” endowed with powerful structures, particularly “fiber products” defining “base change”:

$$\begin{array}{ccc} S' \times_S X & = & X' \dashrightarrow X \\ \downarrow & & \downarrow \\ S' & \longrightarrow & S \end{array}$$

The fascinating notion of continuous family of geometric objects

- Before I discovered the theory of schemes, I only knew the notion of continuous map
 $S \longrightarrow X$ between topological spaces S, X ,
which could be considered as a notion of
continuous family of elements of X
indexed by elements of S .
- Then I discovered that, in scheme theory, a scheme morphism

$$\begin{array}{c} X \\ \downarrow \\ S \end{array}$$

- which is both “flat” and “proper”
(two perfectly defined properties)
can be considered as a continuous family of
algebraic varieties
indexed by the points of S .
- Even more incredibly, if $S = \operatorname{Spec}(\mathbb{Z})$,
points of S were prime numbers $2, 3, 5, 7, \dots$

Discovering “Récoltes et Semailles” after EGA and SGA

- While I had studied EGA and SGA for several years and was preparing a PhD, I began reading “Récoltes et Semailles”:
“Harvests and Sowings: Reflections and Testimonials on a Mathematician’s Past”.
- I discovered that Grothendieck was as singular as a man as his mathematics were singular.
- It seemed that nobody had loved mathematics as much as him, but he was also a writer, a poet and a thinker, and his love and passionate search for truth went far beyond mathematics.
- For him, mathematics consisted much more in definitions than in proofs.
- He was accusing well-known mathematicians and the “mathematical community” as a group.
- He was defending “orphans”: “orphan” theories and “orphan” people.

A most important example of orphan theory

(which has become mainstream 40 years later . . .)

The “six operations” formalism:

This is Grothendieck’s general theory of cohomology
and cohomological duality

(which generalizes simultaneously

Poincaré duality in topology

and

Serre duality in geometry).

Ordinary inner operations:

$\otimes_{\mathcal{O}}, \mathcal{H}om_{\mathcal{O}}$ and their “derivations” $\otimes_{\mathcal{O}}^L, R\mathcal{H}om_{\mathcal{O}}$
well-defined over any topos $\mathcal{E}$
endowed with an inner commutative ring $\mathcal{O}$.

Ordinary outer operations:

f^*, f_* and their “derivations” Lf^*, Rf_*
well-defined for any topos morphism $\mathcal{E}' \xrightarrow{f} \mathcal{E}$
complete with a morphism $f^* \mathcal{O} \rightarrow \mathcal{O}'$
relating inner commutative rings $\mathcal{O}, \mathcal{O}'$ on $\mathcal{E}, \mathcal{E}'$.

The “extraordinary operations” $f_!, f^!$

- They are associated to some topos morphisms (not all)

$$\mathcal{E}' \xrightarrow{f} \mathcal{E}$$

completed with a morphism $f^* \mathcal{O} \rightarrow \mathcal{O}'$.

- Mysteriously, they are usually well-defined in the context of “geometric categories” coming from classical or natural mathematical practice.
- They exist only at the level of “derived categories”:

$$D(\mathcal{M}od_{\mathcal{O}'}) \begin{matrix} \xrightarrow{f_!} \\ \xleftarrow{f^!} \end{matrix} D(\mathcal{M}od_{\mathcal{O}})$$

One of the deepest drawings of mathematics:

$$\begin{array}{ccc} D(\mathcal{M}od_{\mathcal{O}'})^{\text{op}} & \xrightarrow{R\mathcal{H}om_{\mathcal{O}'}(\bullet, f^! \mathcal{L})} & D(\mathcal{M}od_{\mathcal{O}'}) \\ \downarrow f_! & & \downarrow Rf_* \\ D(\mathcal{M}od_{\mathcal{O}})^{\text{op}} & \xrightarrow{R\mathcal{H}om_{\mathcal{O}}(\bullet, \mathcal{L})} & D(\mathcal{M}od_{\mathcal{O}}) \end{array}$$

This is “global duality” ...

An example of orphan person: Olivier Leroy

- He was a student in Montpellier of Carlos Couton-Carrère (a former student of Grothendieck).
- He worked on toposes.
- Grothendieck found him extraordinary talented but he never got any university position and eventually passed away at early age.
- After his death, some mathematical manuscripts were found in his belongings.
- In one of his manuscripts, he had shown that
the notion of subtopos
solves the famous
Banach-Tarski paradox.

Banach-Tarski paradox and Leroy's theorem

Banach-Tarski paradox. –

For a measure μ of open subsets U of a space X
which is invariant by a symmetry group G ,
it is not possible in general to extend μ to all subsets $X' \hookrightarrow X$
in a way which respects both G -invariance and additivity

$$\mu(X'_1) + \mu(X'_2) = \mu(X'_1 \cup X'_2) + \mu(X'_1 \cap X'_2).$$

Leroy's solution. –

There exists a natural extension of μ to all subtoposes

$$\mathcal{E}' \hookrightarrow \mathcal{E}_X \quad \text{of} \quad \mathcal{E}_X = \text{topos defined by } X$$

which respects both G -invariance and additivity.

This measure restricts in particular to subsets $X' \hookrightarrow X$
as any subset $X' \hookrightarrow X$ defines a subtopos

$$\mathcal{E}_{X'} \hookrightarrow \mathcal{E}_X.$$

Explanation:

In general, one has $\mathcal{E}_{X'_1 \cap X'_2} \neq \mathcal{E}_{X'_1} \cap \mathcal{E}_{X'_2}$.

It may even happen that

$$X'_1 \cap X'_2 = \emptyset \quad \text{but} \quad \mu(\mathcal{E}_{X'_1} \cap \mathcal{E}_{X'_2}) > 0.$$

From Grothendieck's geometry to the Langlands program

- Beginning with my PhD thesis under Prof. Gérard Laumon, I made use of Grothendieck's powerful theories, following Drinfeld, to ultimately prove Langland's correspondence for groups GL_r over function fields F_C of smooth projective curves C over finite fields $\mathbb{F}_q$.

Main ingredients:

- Consider the classification problem of Drinfeld's rank r "shtukas" over such a curve C .
- Show this classification problem has a solution in Grothendieck's sense, as a "classifying" algebraic stack Sht_C^r .
- Consider the cohomology spaces (with ℓ -adic coefficients) of the "étale toposes" associated to these algebraic stacks Sht_C^r , with the structural double actions on them of Galois groups $\mathcal{G}al(F_C)$ and linear arithmetic groups $GL_r(\mathbb{A}_{F_C})$.
- Use tools from Grothendieck's cohomological duality theory, especially the "Grothendieck-Lefschetz fixed points formula", to show that the ℓ -adic cohomology spaces of Sht_C^r endowed with the natural actions of $\mathcal{G}al(F_C)$ and $GL_r(\mathbb{A}_{F_C})$ realize Langlands' correspondence.

What is a “classification problem” in Grothendieck’s sense?

Studying Grothendieck, I discovered that, for him,
a well-posed classification problem
consists in a definition of a general notion of
“continuous family of structures of some type $\mathbb{T}$
indexed by points of spaces S of some type $\mathcal{C}$ ”
such that:

- (1) For any space S of the type $\mathcal{C}$,
continuous families of structures of type $\mathbb{T}$ induced by S
make up a set
 $M_{\mathbb{T}}(S)$ (or more generally a “category” $\mathcal{M}_{\mathbb{T}}(S)$).
- (2) For any map between spaces S, S' of type $\mathcal{C}$
 $S' \xrightarrow{s} S$ which respects their geometric structures,
there exists an associated “change of parametrization” map
 $s^* : M_{\mathbb{T}}(S) \longrightarrow M_{\mathbb{T}}(S')$
(or more generally a functor $s^* : \mathcal{M}_{\mathbb{T}}(S) \rightarrow \mathcal{M}_{\mathbb{T}}(S')$).

What is a “classifying space” in Grothendieck’s sense?

Definition. –

Consider a classification problem in Grothendieck’s sense

$$\begin{array}{ccc} \mathcal{C} & \longrightarrow & \text{Set (or Cat),} \\ \left. \begin{array}{c} \text{space } S \\ \text{of } \mathcal{C} \end{array} \right\} & \longmapsto & \left\{ \begin{array}{l} M_{\mathbb{T}}(S) = \underline{\text{set}}, \text{ or} \\ \mathcal{M}_{\mathbb{T}}(S) = \underline{\text{category}}, \end{array} \right. \\ \left. \begin{array}{c} \text{geometric map } S' \xrightarrow{s} S \\ \text{of } \mathcal{C} \end{array} \right\} & \longmapsto & \left\{ \begin{array}{l} M_{\mathbb{T}}(S) \xrightarrow{s^*} M_{\mathbb{T}}(S'), \text{ or} \\ \mathcal{M}_{\mathbb{T}}(S) \xrightarrow{s^*} \mathcal{M}_{\mathbb{T}}(S'). \end{array} \right. \end{array}$$

Then a (necessarily unique) “classifying space” for this classification problem

is a space $S_{\mathbb{T}}$ in $\mathcal{C}$

endowed with a so-called

universal continuous family indexed by $S_{\mathbb{T}}$

$U_{\mathbb{T}}$ in $M_{\mathbb{T}}(S_{\mathbb{T}})$ or $\mathcal{M}_{\mathbb{T}}(S_{\mathbb{T}})$,

such that, for any space S of $\mathcal{C}$, the map (or functor)

$$\begin{array}{ccc} \left\{ \begin{array}{c} \text{geometric maps} \\ S \rightarrow S_{\mathbb{T}} \text{ in } \mathcal{C} \end{array} \right\} & \longrightarrow & M_{\mathbb{T}}(S) \text{ or } \mathcal{M}_{\mathbb{T}}(S), \\ (S \xrightarrow{s} S_{\mathbb{T}}) & \longmapsto & s^* U_{\mathbb{T}} \end{array}$$

is one-to-one (or is an equivalence).

Classifying toposes

In 2010, I learnt from a young post-doc, Olivia Caramello, the following theorem which had been proved in the mid 1970's by Makkai and Reyes, building on previous work of Grothendieck, Monique Hakim, Joyal, ...

Theorem. –

Consider an arbitrary “first-order theory” $\mathbb{T}$ presented in “geometric form”.

Then:

- (i) For any “topos” $\mathcal{E}$ (defined by Grothendieck as the most general notion of space),
continuous families of structures of type $\mathbb{T}$ indexed by $\mathcal{E}$
make up a category $\mathcal{M}_{\mathbb{T}}(\mathcal{E})$.
- (ii) Any topos “morphism” $\mathcal{E}' \xrightarrow{f} \mathcal{E}$
(defined by Grothendieck as the natural generalization of continuous maps)
defines a functor of “change of parametrization” of continuous families
 $f^* : \mathcal{M}_{\mathbb{T}}(\mathcal{E}) \longrightarrow \mathcal{M}_{\mathbb{T}}(\mathcal{E}')$.
- (iii) Whatever $\mathbb{T}$, this classification problem defines
a “classifying topos” $\mathcal{E}_{\mathbb{T}}$
which is uniquely determined up to equivalence.

Remark:

For any such a theory $\mathbb{T}$, i.e. any formal language (of “syntactic” nature),
 $\mathcal{E}_{\mathbb{T}}$ incarnates its “semantics” as a generalized space.

Caramello's technique of "toposes as bridges"

Basic facts. –

- (i) Following Grothendieck, any "geometric datum" consisting in
- a category $\mathcal{C}$, i.e. a world of spaces of some type,
 - a "topology" (in a generalized sense) J on $\mathcal{C}$
- defines a topos $\mathcal{E} = \hat{\mathcal{C}}_J$.
Conversely, any topos $\mathcal{E}$ is associated to
infinitely diverse geometric data $(\mathcal{C}, J)$.
- (ii) Any formal language $\mathbb{T}$ of "first-order geometric" type
defines a "classifying topos" which incarnates its semantics.
Conversely, any topos $\mathcal{E}$ is equivalent to the classifying topos $\mathcal{E}_{\mathbb{T}}$ of
infinitely diverse formal languages $\mathbb{T}$.

Principles for the "bridge technique". –

- (1) Consider invariants of toposes.
- (2) Express them or compute them in terms
of different types of presentation by geometric data $(\mathcal{C}, J)$
or by formal languages $\mathbb{T}$.
- (3) Combine these expressions or computations
with equivalences of toposes associated to concrete situations.

An example of the “bridge technique”

Basic geometric fact. –

Grothendieck has defined a notion of “subtopos” of a topos $\mathcal{E}$

$$\mathcal{E}' \hookrightarrow \mathcal{E}$$

which generalizes the notion of subspace of a space.

Theorem (Grothendieck). –

For any presentation of a topos $\mathcal{E}$ by a geometric datum $(\mathcal{C}, J)$

$$\mathcal{E} \cong \widehat{\mathcal{C}}_J,$$

subtoposes $\mathcal{E}' \hookrightarrow \mathcal{E}$

have a one-to-one correspondence with

topologies J' on the underlying category $\mathcal{C}$ which refine J .

Theorem (Caramello). –

For any presentation of a topos $\mathcal{E}$ as classifying a theory $\mathbb{T}$

$$\mathcal{E} \cong \mathcal{E}_{\mathbb{T}},$$

subtoposes $\mathcal{E}' \hookrightarrow \mathcal{E}$

have a one-to-one correspondence with

{ systems of extra axioms added to $\mathbb{T}$,
considered up to provable equivalence. }

From Grothendieck's geometry to AI following toposes

Application to provability problems. – *The double equivalence*

$$\left\{ \begin{array}{c} \text{topologies} \\ \text{on categories} \end{array} \right\} \leftrightarrow \left\{ \begin{array}{c} \text{subtoposes} \\ \text{of a topos} \end{array} \right\} \leftrightarrow \left\{ \begin{array}{c} \text{systems of extra axioms,} \\ \text{considered up to provable equivalence} \end{array} \right\}$$

is constructive and allows computation.

It allows to transform any provability problem in math
into problems of generating topologies.

Application to “syntactic learning”. –

The definition and construction of classifying toposes $\mathcal{E}_{\mathbb{T}}$
relates in a perfectly precise and rigorous way

- the problem of describing and representing
any type of data we are interested in,
- the problem of classifying data of any type.

Application to the possible creation of a theory of semantic information. –

Such a theory could and should be based on the construction

$$\mathbb{T} \longmapsto \mathcal{E}_{\mathbb{T}}.$$

Application to “meta-learning”. – *The notion of relative topos*

$$\mathcal{E}' \rightarrow \mathcal{E}$$

provides a mathematical model for “meta-learning”,

i.e. extracting information at higher levels of abstraction.

Helping to create and develop the “Grothendieck Institute”

- The “Grothendieck Institute”

is a Scientific Foundation.

It was created in 2022 by Prof. Olivia Caramello
and is recognized by the Italian government.

Purposes of the Grothendieck Institute. –

- (1) To disseminate the knowledge of Grothendieck’s

- published work,
- manuscripts,
- thought.

- (2) To contribute to the continuous development of

- mathematics inherited from Grothendieck,
- general unifying theories such as those he introduced,
- in particular Topos Theory and its applications.

Remark:

Grothendieck’s thought goes beyond mathematics.

For instance, his book

“La Clef des Songes” (“The Key of Dreams”)

has deep philosophical content.